Design and Implementation of Highly Scalable E-mail Systems

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Apologies

• Less “Implementation”
• More “Fundamentals & Architecture”
  – This stuff is hard
  – This stuff is surprisingly hard, even for experienced professionals
• Nick unable to attend

I know now just how much hubris it really took to choose a title like this.

When I submitted my proposal for this talk, Skynet was just beginning the design phase for their “next generation” mail system, and I was quite confident that by this time, I would simply be talking about history — the “been there, done that, took some pictures, here’s my album if you’d care to look at it” sort of thing.

Boy, was I wrong. In fact, Skynet is just now starting the implementation phase for the next-generation mail system (based on the architecture presented in this talk), and no one knows yet just how everything is going to turn out.

Of course, everyone talks about their having an IMAP-based mail system that will scale to millions of users, but no one actually has such a beast. The largest I am personally familiar with is about 200,000 users. When Skynet turns their system on, it should immediately be about one million users, since that’s about what we’ve got across all our various mail systems today, and this is intended to take over from all of them.
Outline

- Review of Major Publications
- Review of Typical POP3 Implementations
  - Enhancements
- Contrast with IMAP
  - Implications of protocol differences
- Functional Architecture
- Detailed Architecture
Information Sources

- Academia
  - Build vs. Buy
    - Frequently re-invent the wheel
  - Small Scale
  - Occasionally revolutionary

- Commercial
  - Buy vs. Build
    - Time-to-market crucial
  - Large Scale
  - Usually Evolutionary
  - Any revolutions are usually in the area of scaling
The purpose of this paper is to help you develop an appreciation for what has happened so far in this industry, and outline an architecture that will help fill in the space shown by the red circle.
In my opinion, before you can lay new ground, you need to fully understand everything that has come before. Therefore, we start off with a review of notable papers, publications, and presentations that we have been able to find so far.

In the interest of brevity, we will actually skip everything on the left-hand side of this chart. I encourage you to review all this material yourself (the entries can be found in the Bibliography), so that you can better appreciate what all it takes to do something like this correctly.

However, I do still have the forty-something slides that comprised this material, and plan on making them available via my web site at <http://www.shub-internet.org/brad/papers/>.
Publications Review

• POP3 Mail Systems
  – Grubb96
  – Christenson97
  – Horman99
Review: Grubb96

- Problem
  - NFS mail spool/hub configuration using 7th edition mailbox (mbox) format for ~5000 users could not scale to ~20,000 users
Review: Grubb96

- Solutions
  - Front-end MXes handle incoming communications
  - Back-end servers handle mailboxes
    - Front-ends “trickle” feed via smaller number of cached connections to back-end servers
  - Separate syslog data onto separate disk
  - Tweak kernel, NFS server, & NFS client settings
  - Change client config to use mailhub name based on userid via DNS CNAME records
Review: Grubb96

- Solutions, continued
  - Implement additional mailhubs to serve chunks of user community based on CNAMEs
  - Turn on POP3 & IMAP2bis w/ 7th edition mailbox (mbox) format on each new “post office” server
  - Provide users with POP3/IMAP clients
  - Turn off NFS
  - Convert POP3/IMAP2 mbox → POP3/IMAP4 Cyrus on each “post office” server
Review: Grubb96

- Applicability
  - Does cover entire mail system, not just MTA
  - Doesn’t really tell us anything about how POP3/IMAP system is managed
  - Doesn’t scale
    - Users have to know too much about post office configuration
    - Requires CNAME RR for each customer
  - Mixes inbound and outbound services on same machines
Review: Christenson97

• Problem
  – Existing information on architecture for robust large-scale mail systems is scarce, doesn’t address key issues, and doesn’t scale to required levels
Review: Christenson97

- **Solutions**
  - Front-end MXes handle external communications
    - Secondary MXes do not attempt delivery to back-end, in case there is a problem with deliveries
    - Front-end MXes do not authenticate recipient names
  - **All** machines are dataless
  - Modify LDA to handle authentication methods, mailbox formats and quotas
Review: Christenson97

- Solutions, continued
  - Back-end servers do accept outgoing SMTP mail
    - Do not do local delivery, pass to inbound MXes instead
    - POP3 code must also be modified to know about authentication methods and mailbox format
  - All data (mailboxes and sendmail mqueues) stored on NetApp NFS servers
  - Mail spool directories hashed and split across multiple NFS servers

Note that this is a statement of the original state of the network.

Later, management wanted to increase the benchmark performance of this system, and the sendmail mqueues were moved to solid-state disks that were attached locally to the respective machines.

This didn’t measurably increase the performance from the user perspective, but did significantly improve the performance of the system on certain benchmarks, which was important for certain magazines.
Review: Christenson97

- Solutions, continued
  - Dynamically balance mailboxes or expand capacity
    - Both POP3 daemon and LDA know about “old” vs. “current” mailbox location
    - POP3 daemon moves mailbox if necessary
  - POP3 & LDA modified to use database for user authentication, avoiding use of `/etc/passwd`
  - Cluster & fail-over for user authentication database
Review: Christenson97

• Solutions, continued
  – NFS file locking doesn’t work reliably
    • Replace w/ lockfiles on separate shared NFS server
      – Uses semantics of open() system call with exclusive write
    – Lock system needs to be replaced to scale further
      • Custom clustered servers w/ shared RAID & unbuffered writes
      • Query different lock servers for different ranges of mailbox names
Review: Christenson97

- Applicability
  - Does cover entire mail system in centralized fashion
  - NFS servers are SPOFs
  - UDP & RPC are major security hazards
  - Customized code is expensive to maintain
  - Specific to POP3, does not cover IMAP

Obviously, my co-author doesn’t necessarily agree with these points, since this was the mail system he had built, and is the basis for some mail systems which have since been installed by Sendmail, Inc. for various customers.

I would like to keep an open mind on this issue, and I have a great deal of respect for his experience and expertise, but my own personal experience so far with NFS servers has not been as rosy, and I am much more strongly disinclined to use them in this role.

My co-author is completely correct when he notes that the security issues with UDP & RPC can be resolved by putting the back-end mail servers and the NFS servers that they talk to on a private (unroutable) network, but that still doesn’t eliminate my natural distaste for UDP & RPC.
Review: Horman99

• Goal
  – Define architecture to scale mail systems transparently to multiple servers
Review: Horman99

• Solutions
  – Multiplex SMTP
    • Single layer
      – If recipient not local, must forward to correct server
      – With growth, amount of forwarding approaches 100%
    • Dual layer
      – No local recipients on front-end servers
      – Must always forward to correct back-end server
  • Add layer 4 load-balancing switches to hide number of machines accepting SMTP connections
Review: Horman99

• Solutions, continued
  – Multiplex POP3 & IMAP
    • Single layer
      – Must handle local connections
      – Must also proxy for remote connections
    • Dual layer
      – Dedicated content-free proxies
Review: Horman99

• Solutions, continued
  – Mailbox migration
    • Calculate metric for each server over reasonable time
    • Migrate only if a server deviates significantly from avg.
    • Order users by decayed metric cost
      – How long are migrations remembered?
      – How long since this mailbox migrated?
    • Generate user list probabilistically
Review: Horman99

- Solutions, continued
  - Mailbox migration, continued
    - Move from most heavily loaded server to least heavily loaded server(s)
    - Move only if result would not push recipient over average
    - Continue with next most heavily loaded server until no more migrations are possible
Review: Horman99

- Applicability
  - Covers only POP3 and not IMAP
  - Proper load balancing requires programming for peaks, not long-term averages
  - Focuses exclusively on “free” or “cheap” solutions
  - Too much time/space spent on less important issues
  - Not enough detail provided where needed
Publications Review

- IMAP Mail Systems
  - Stevens97
  - Beattie99
Review: Stevens97

- Statistics
  - 60,000 accounts
  - 4,000 peak concurrent logins
  - 1.4 million logins per month
  - 500,000 messages/day
  - 1,083 peak messages/minute
    - 65,000 peak messages/hour

I looked long and hard at this paper, and just didn’t see a whole lot in the way of useful technical information that could be analyzed. However, I did find these statistics to be of interest.
Review: Beattie99

- Goals
  - Implement and document replacement mail system for ~30,000 users
    - Reliable
    - Secure
    - IMAP & SMTP
    - Web interface available
    - Quotas
Review: Beattie99

• Solutions
  – Mix & match software on cluster of commodity computers running Unix-like OS
    – UW imapd
    – Exim
    – Apache/mod_perl
    – WING (Web IMAP/NNTP Gateway)
    – PostgreSQL
    – BIND
    – Custom account & cluster management tools
Review: Beattie99

• Solutions, continued
  – Two front-end servers are firewalls & nameservers
    • Configured for fail-over
  – IMAP servers hold all per-user filestore
    • IMAP, POP3, & SMTP (public)
    • NFS export to other nodes (private)
      – Vacation messages
      – Forward files
      – Personal home page links
Review: Beattie99

- Solutions, continued
  - WING servers hold only temporary data
    - HTTP (public)
    - IMAP & NFS to IMAP/NFS servers (private)
    - SQL to front-end/firewall servers (private)
  - Each user has DNS entry
    - username.herald.ox.ac.uk
    - CNAME alias to home IMAP node
Review: Beattie99

- Solutions, continued
  - Front-end machines are
    - Cluster nameservers
    - SMTP & HTTP login gateways
    - DBMS servers for all user config data
    - Generate mailer tables and push to other nodes
Review: Beattie99

- Solutions, continued
  - Security
    - Front-ends are firewalls
    - IMAP & WING servers trust front-ends 100%
    - IMAP servers export ~foo/wing directory owned by httpd for each user foo
      - Automap games to handle mounts
    - Break-in on WING servers allows modification of forward files, vacation messages, & personal links but NOT mail
Review: Beattie99

• Solutions, continued
  – Failure analysis
    • IMAP
      – Mail stored on RAID5
        » Immune to single disk failure
        » If node dies, all users on that node lose access
    • WING
      – Current sessions die
      – 1/n login attempts fail until server manually removed from lists
    • Switch = SPOF
Review: Beattie99

- Solutions, continued
  - Failure analysis
    - Front-end
      - DNS continues
      - IP traffic dropped but can reconnect
      - SQL failover currently manual
        » Lose config changes since last sync
  - Changes
    - Added outbound mail relay servers to speed up acceptance of mail from dumb clients
Review: Beattie99

- Statistics
  - Recent average week
    - 2 IMAP servers, 2 WING servers
    - 82,000 total connections to IMAP servers
    - 113,000 mail deliveries by IMAP servers
      - 95,000 local
      - 18,000 outgoing
    - 26,000 outgoing messages from WING
    - 66,000 IMAP sessions (including 38,000 WING)
    - 120,000 POP3 sessions
Review: Beattie99

- Applicability
  - Very small scale
    - We have \(~7.5x\) their # of users
    - We do \(~38x\) their number of inbound mail messages
    - We do \(~35x\) their number of local mail deliveries
    - We do \(~64x\) their number of outbound mail messages
    - We don’t know how many more POP3 sessions we do
      - Too expensive to track
Review: Beattie99

- Applicability
  - Not scalable, not enough functional decomposition of services
    - Front-end/firewall/nameserver/user meta-data server doing way too much
    - IMAP servers should not be used as outbound mail relays
    - IMAP servers should not be used as NFS servers
Publications Review

- Distributed
  - Yasushi99
Review: Yasushi99

• Goals
  – Build and describe distributed, replicated, clustered, automatically load-balanced, functionally homogenous mail system
Review: Yasushi99

• Solutions
  – Use commodity hardware and OS
  – Write all custom application code
  – Mailboxes fragmented at message level
    • Replicated across two servers
    • Distributed across as many as four servers
  – All servers run all protocols
    • SMTP in & out, POP3, IMAP, User metadata database
Review: Yasushi99

- Solutions, continued
  - Soft limit of four distributed servers can be exceeded if one or more nodes is down
  - Some affinity of distributed servers is maintained to reduce latency
  - Automatically discover new resources
  - Detect and route around failures automatically
  - Balance cluster automatically across all nodes
Review: Yasushi99

• Solutions, continued
  – Claims to be lock-free because POP3 and IMAP require only convergence to consistency over time
  – “Load” defined as boolean + integer
    • Disk full or not?
    • Total number of outstanding potential I/O requests
  – Node with full disk is always considered to be “very loaded”
    • Used only for reading and deleting mail
Review: Yasushi99

- Solutions, continued
  - Testing methodology
    - Avg. msg size 4.7KB w/ fat tail to 1MB
    - SMTP traffic = 90% of load
    - POP3 traffic = 10% of load
    - Compare against sendmail 8.9.3 + ids-popd-0.23
    - Custom load-generation programs
    - POP3 test program collects and deletes all mail for user
    - Linux async writes are used
Review: Yasushi99

- Solutions, continued
  - Testing results
    - One node w/ no replication and one IDE disk could handle ~23 msgs/sec.
    - Adding two SCSI disks to single node, it could handle ~105 msgs/sec.
    - Two nodes w/ one IDE and two SCSI disks each could handle ~38 msgs/sec. w/ replication, ~48 msgs/sec. w/ simulated NVRAM for coordinator log
Review: Yasushi99

• Solutions, continued
  – Testing implications
    • @ ~105 msgs/sec. per node, ~62 nodes could saturate
      1Gbps network, w/ ~562 million msgs/day
      – ~6500 msgs/sec. aggregate
    • With replication, this drops to ~5200 msgs/sec. aggregate,
      and ~450 million msgs/day on ~108 NVRAM nodes or
      ~137 non-NVRAM nodes
Review: Yasushi99

• Applicability
  – Throws out all previous application work
    • 100% new, untrusted code
  – Can’t list 100 IP addresses in DNS for POP services
    • Won’t fit into 512 byte UDP packets
  – Can’t list 100 IP addresses in DNS for MX services
  – Forced to use proxy front-ends or L4 load-balancing switches to hide the number of servers
Review: Yasushi99

- Applicability
  - Microsoft OSes only ever use the first IP address, then cache forever (until reboot)
  - Forced to use L4 load balancing switches
    - Must be set up in HA/failover mode
    - May have application proxies behind them
  - Some SMTP MTA or resolver implementations are equally dain-bramaged
    - L4 load-balancing switches in front of MXes
Review: Yasushi99

- Applicability
  - Can’t get around DNS UDP packet size restrictions with multiple IP addresses per name
    - If connection refused, skip to next name
    - Iff connection timed-out, go to next IP address for same name
  - At ~2 min. TCP timeout per IP address, 45 IP addresses = 90 minutes to timeout
    - If you have a queue runner fired off every 60 minutes, you ultimately wind up with all memory taken up and no mail flow
Review: Yasushi99

- Applicability
  - Did not use standard benchmarking tools
    - May or may not be valid to create own tools, but needs justification
  - Fundamentally, locking IS required
    - Users simply will not accept messages appearing and disappearing and reappearing again
    - Requires serialization which violates most basic principles espoused
Review: Yasushi99

- Applicability
  - Did not test suitable array of MTAs, POP3 daemons, message size and arrival distributions, mailbox sizes, etc...
    - Did not even prove special case, much less general case
    - Anybody can select bad special case and demonstrate superiority
    - To claim general superiority, you must test across a much broader array of variables
Review: Yasushi99

- Applicability
  - IMAP implementation is only a subset — does not include shared folders
    - Perhaps possible in small academic environment
    - Simply not acceptable in large commercial environment
  - SMTP server holds sender open while all writes are completed
    - Violation of RFC 1123, section 5.3.2?
    - All other MTAs accept first, then deliver in background
Review: Yasushi99

• Applicability
  – Each server must implement all protocols
    • Doesn’t allow for scaling of each part independently
  – Load discovery protocol is broadcast-based
  – Uses Linux async writes
    • Violation of RFC 1123, section 5.3.3
    • Replication already used to address lower reliability of commodity hardware, OS, and custom application code
Review: Yasushi99

- Applicability
  - Peak sustained rates do not scale linearly
    - msgs/sec. \rightarrow msgs/min. \rightarrow msgs/hr. \rightarrow msgs/day
      - msgs/hr. \times 10 \approx msgs/day
Review: Yasushi99

- Applicability
  - Good things
    - Splitting mailboxes at message level
    - Replicate messages to at least two servers
    - Distribute messages across up to four servers
    - Dynamically distribute messages to least loaded servers
    - Calculate “load” based primarily on current and potential disk I/O operations

While I believe that this paper has many fundamental flaws, it still was the most interesting that I’ve seen in a long time.
Skynet Statistics

- POP3 Mail Server
  - 285,000 Accounts
  - 225,000 Mailbox files
  - 600,000 Aliases
  - 6800 Domains
  - 150 GB Total mailbox storage
    - 1 GB Overhead
Skynet Statistics

- POP3 Mailbox Sizes
  - 80,000 Empty
  - 690 KB Average
  - 9282 bytes Median (50th percentile)
  - 1.1 MB 90th percentile
  - 3.35 MB 95th percentile
  - 12 MB 99th percentile
  - 42.1 MB 99.9th percentile
Skynet Statistics

- POP3 Connections
  - 100 peak connections/attempts per second
  - 2300 peak connections/attempts per minute
  - 105,000 peak connections/attempts per hour
  - ??? peak connections per day?
  - 13.14 second typical daily average connection time
  - 300 Max total simultaneous connections allowed
Skynet Statistics

Millisecond response times (14 day sample)

<table>
<thead>
<tr>
<th>Protocol</th>
<th>Min</th>
<th>Avg</th>
<th>Max</th>
</tr>
</thead>
<tbody>
<tr>
<td>SMTP</td>
<td>33</td>
<td>672</td>
<td>3600</td>
</tr>
<tr>
<td>POP3</td>
<td>28</td>
<td>185</td>
<td>949</td>
</tr>
</tbody>
</table>

In this case, having high latency for handling of SMTP on this server is okay — the only machine(s) that ever contact it via SMTP are our own front-end mail servers, and we can safely design this part of the system for maximum throughput as opposed to minimum latency.

Please note that this is not the normal case for servers which speak SMTP to external machines.
Skynet Statistics

• Typical messages per day
  – 450,000 inbound SMTP
    • 450,000 POP3 mailbox deliveries
    • 200,000 webmail/freemail
    • 40,000 business SMTP
  – 400,000 outbound SMTP
Skynet Statistics

- Peak messages per hour
  - 48,000 inbound SMTP
  - 42,000 outbound SMTP
Skynet Statistics

• Typical message volume per day
  – 48 GB inbound
    • 25 GB POP3
    • 18 GB webmail
    • 4.5 GB business
  – 48 GB outbound
Skynet Statistics

- Average message sizes
  - 110 KB inbound
    - 60 KB POP3
    - 100 KB webmail
    - 120 KB business
  - 120 KB outbound
Protocol Implementation Analysis

• POP3
  – Typical implementation
  – Qpopper “Server Mode”
  – Indexed Mailbox
  – Login Frequency Limitation
  – Mailbox Directory

• IMAP Differences & Implications
Analysis: Typical POP3

- User login
- Lock mailbox
- Create temp file
- Copy mailbox to temp file
- Truncate mailbox
- Unlock mailbox
- Operate on temp file
  - New messages may come in to mailbox
Analysis: Typical POP3

- User logout
- If any messages are being retained
  - Re-lock mailbox
  - If mailbox not empty
    - Append new messages to temp file
    - Truncate mailbox
  - Merge retained temp file contents onto mailbox
  - Unlock mailbox
- Delete temp file

Does this sound like a lot of operations to you? Well, many of those operations are synchronous meta-data operations, and therefore exceedingly expensive to incur on a mail server.
Analysis: Qpopper “Server Mode”

- User login
- Lock mailbox
- Operate on mailbox
  - New mail messages wait to be added to mailbox
- User logout
Analysis: Qpopper “Server Mode”

- Are messages being retained?
  - Yes
    - Create temp file
    - Merge retained contents of mailbox onto temp file
    - Move temp file to mailbox
  - No
    - Truncate mailbox
- Unlock mailbox

This does require that incoming mail messages may be temporarily delayed a bit while the user is processing their mailbox, but overall I still believe that this is a significant win.
Analysis: Qpopper “Server Mode”

• Improvements
  – Big “win” if no mail is left on server
    • Virtually all synchronous meta-data operations eliminated
  – No “loss” if mail is left on server

• Issues
  – Still have to scan entire mailbox every time user logs in, even if only to tell them they don’t have any new messages
Analysis: Indexed Mailbox

- User login
- Lock index
- Stat index & mailbox
- If index newer, all questions can be answered from index
  - Only need to lock mailbox if messages are deleted
Analysis: Indexed Mailbox

- If mailbox newer
  - Lock mailbox
  - `lseek()` to last position specified by index, then scan and update index
- Otherwise, like Qpopper “Server Mode”
Analysis: Indexed Mailbox

- Improvements
  - Each message read from mailbox is handled by lseek() and large-size read()
  - Greatly increases use of read-ahead cache
  - Assumes that LDA appends only
  - Assumes that LDA & POP3 server are only methods of reading or writing mailboxes
Analysis: Indexed Mailbox

- Problem
  - Still have to update mailbox if messages are retained and message status changes

- Solution
  - In index, separately store header and body start+offset info
  - Store message status in index
  - Generate message status header info on-the-fly
Analysis: Indexed Mailbox + status

• Results
  – Twice as many read operations
  – Fewer write operations
  – More complex POP3 server
    • Probably a big win for leave-on-server
Analysis: Limiting User Login

• Problem
  – Some clients still login too frequently to check their mail

• Solution
  – Require that at least $X$ minutes elapse before you allow updating of index
  – Tune $X$ for pain threshold of your users

We have some clients that log in as frequently as once every five seconds to check their mail. Obviously, this is just too often. But what is a more reasonable number?

We believe that it is reasonable to insist that users wait at least five minutes between login attempts, and since we don’t really want to refuse their login, we instead choose to be careful what we tell them and how, and in what timeframe.
Analysis: Mailbox Directory

- Some POP3 implementations create a directory that comprises the mailbox, and store one message per file
  - Trades smaller number of larger I/O operations for much larger number of smaller I/O operations
  - Avoids mailbox locking issues
  - Creates message locking issues
Analysis: Mailbox Directory

- Problems
  - The I/O operations it creates in trade are all synchronous meta-data operations
    - The most expensive kind
    - The type we most want to eliminate, reduce, or optimize
  - May need to implement directory hashing within mailbox to avoid excessively large directories
Analysis: Mailbox Directory

• Problems
  – Typically has to scan entire directory tree to build mailbox status
    • Must know size of each message
      – Must `stat()` each file or have file size encoded in file name
    • Must know UIDL value for each message
      – Must open and read each file
  – Can solve these problems by using index
    • Still doesn’t eliminate sync. meta-data updates
Analysis: Mailbox Directory

• Claim
  – More NFS-friendly
  – Avoids mailbox locking
  – Mechanism for creating filenames sufficiently unique to virtually eliminate collisions on files
    • Uses “create w/ exclusive ownership” semantics to detect
Analysis: Mailbox Directory

- Reality
  - Christenson97 shows that 7th edition mailbox (mbox) format can also be made NFS-friendly, using same trick
  - Still have issues with sync. meta-data updates
    - Now problem for NFS server vendor?
  - Does not solve locking problems with message changes, moves, or deletions
  - Mailbox locking not really a problem

In the grand scheme of things, while locking is certainly necessary (and I believe that the POP3 protocol requires mailbox locking), if properly implemented this shouldn’t pose that big of a problem.

Switching to message locking seems excessive and unnecessary for POP3 environments, especially since it creates far more synchronous meta-data operations in turn for what it supposedly gives you, and it’s synchronous meta-data operations that we most want to eliminate or reduce whenever possible.
Implications

- POP3
  - Only one reader process at a time
    - Can safely lock entire mailbox
  - Only one writer process at a time
    - Can safely lock entire mailbox
  - Long-term mail storage is local to user
  - Large sites may not allow “leave on server”
    - Otherwise mitigated by quota or expiration mechanisms
Implications

- IMAP
  - There will be more than one simultaneous reader and/or writer process
    - Cannot lock entire mailbox
    - Must lock at message level or below
  - Long-term mail storage is centralized
    - Only cached locally
Implications

- Solutions
  - Easiest way to deal with message locking is to avoid 7th edition mailbox (mbox) format
  - Use mailbox directory instead, but can use folders
    - One message per file
    - Some typical POP3 enhancements not applicable
  - However, so long as lock mechanism is shared by LDA & IMAP server, can avoid file locking and use database instead

For IMAP, you need message locking anyway, so the use of a mailbox directory format is natural, although I still believe that file locking with a mailbox directory requires an excessive amount of work and should instead be pulled into a proper database (more on this later).
Scaling Growth

• Problem
  – Number of users is increasing
  – Number of messages sent/received per user is increasing
  – Average size of messages is increasing
  – Length of retention of messages increasing
    • Due to centralized storage of mailboxes
Scaling Growth

- Result
  - Disk storage requirements increasing exponentially
  - Number of I/O operations increasing exponentially
Scaling Growth

• Mitigating Factors
  – Disk storage space increasing exponentially

• Complications
  – Disk rotational speed increasing
    • But not increasing very fast
  – Track-to-track latencies improving
    • But not improving very quickly
Scaling Growth

• Result
  – Disk storage requirements still increasing
    • Not quite as bad
  – Number of I/O operations increasing exponentially
    • Our main killer before
    • Will become bigger and bigger bottleneck
Scaling: Future Improvements

- Single Instance Message Store
  - If storing message per file, store message only once per machine and hard link other recipients to same file
    - Reduces I/O bandwidth requirements
    - Doesn’t reduce sync. meta-data updates since linking to an existing inode requires just as much directory update work as creating new file
Scaling: Future Improvements

- Multi-session Single Instance Message Store
  - Generate MD5 or SHA-1 hash of message
  - Already in system?
    - Yes
      - Compare binary files, store if different, link otherwise
    - No
      - Store
  - Further reduces disk storage capacity issues
  - **Increases** synchronous meta-data I/O

Overall, I don’t believe that multi-session single instance message store is a viable technology to be pursued. It reduces things that we don’t care so much about (disk storage requirements), and it increases the things we do care most about (synchronous meta-data I/O).

The point of bringing this up now is to show you that this is an option that some vendors are touting, but to also demonstrate its limitations, so that you can see through their vacuous claims.
Scaling: Future Improvements

• Multi-session Single Instance in Bodypart Store
  – Recursively parse MIME message structure, store bodypart-per-file
    • For attachments, insensitive to trivial changes in body
    • Allows you to replace base64 or quoted-printable with binary
    • Allows you to “invisibly” compress data
    • Further reduces disk storage requirements
    • Still doesn’t address issues of sync. meta-data updates
Scaling: Future Improvements

• Use Database for Everything  
  – Eliminates sync. meta-data I/O problems  
• Problem  
  – No database handles BLOBs properly  
  – Large scale database reliability problems?

For example, how many hosts are there on the Internet? How many of them have databases on them? Of them, how many have large databases on them? Now, sum the overall reliability across those relatively few machines with large databases.

Now, let’s go back to the total number of hosts on the Internet — how many do you think have filesystems on them? Pretty much all of them, right? Now, sum the overall reliability of those filesystems across all those machines.

Now, let’s compare these two numbers. I submit that the overall total reliability of large databases is many orders of magnitude less than the overall total reliability of filesystems.
Scaling: Future Improvements

- Use Message “heap”
  - Use INN timecaf/timehash-style files instead of message-per-file
    - New message comes in
      - Append to one of small number of large files
      - Update database index
    - Message is deleted
      - Mark space as available
      - Reclaim empty space at time of reduced load
Scaling: Future Improvements

- Message “heap”, continued
  - Virtually eliminates all sync. meta-data updates
  - Could potentially be combined with previous single-instance-store ideas
    - Probably not worth it
  - Does increase maintenance overhead

Of course, now you’ve just thrown out all that previous experience with well-tested filesystems code.

However, since you don’t need everything that a full filesystem comprises, and you don’t need everything that a full-blown RDBMS provides, you can pick and choose the parts that work from both camps, and you should still be able to put something together that will work reasonably well.

After all, there are plenty of INN and Diablo news servers out there that are handling a full 200GB+ full feed these days with similar technologies, and they seem to work pretty well.

So long as you build in enough redundancy that you can recover the database portion should it get corrupted, it seems to me that this would be a good balance between storing everything message-per-file in a filesystem, and storing everything as BLOBs in a database.
Scaling: Future Improvements

- From Yasushi99
  - Break mailboxes into component messages
    - Replicate messages to at least two servers
    - Distribute messages across four or fewer servers
  - Doesn’t help address either disk storage or sync.
    - Meta-data issues
  - Does address issues of reliability, load-balancing,
    - Speed, and perceived quality of service

My co-author believes that once you start mirroring messages or mailboxes to separate servers, you must start doing integrity checks on every read operation, and get into expensive “voting” schemes, etc….

My belief is that this is not an integrity issue but more like doing RAID 1 (mirroring) at the filesystem level, and that you can do the same sort of thing — round robin your read operations, or apply them to the least heavily loaded server, etc….
Scaling: Future Improvements

- Yasushi99, continued
  - Could be combined with INN timecaf/timehash-like message “heap”
  - Could calculate “load” for re-balancing of messages on different criteria
    - Old messages could be migrated to specialized servers with more disk space, perhaps less disk I/O capacity
Best Current Practice

• Per message store server
  – Single instance message store
    • Hard links for multiple recipients of same message
  – Hashed mailbox directories
    • Two base-32 chars per subdir = 1024 max per dir
      – Minimizes path length
  – Message locks in fast and reliable database
    • Berkeley db, not SQL
Best Current Practice

- Per message store server, continued
  - Most important headers and MIME structure in database
    - Most meta-data queries answerable from database
  - User mailbox on single server (cluster)
  - Archive all messages at appl. level, if req’d
  - Clustered servers for HA

Mailbox mirroring may also become a “Best Current Practice”, if we can resolve the issue of whether an integrity check needs to be performed on each read operation, etc.…
Best Current Practice

- User meta-data database kept outside of message store servers
- Minimize interface protocols
- Use application proxies to distribute traffic across $n$ number of message store servers
- Use Layer 4 load-balancing switches in HA mode to hide number of application proxies
Best Current Practice

- Everything becomes LEGO™ building blocks
- However, scaling is still not quite linear
  - 1 million users = one servers
  - 10 million users != ten servers
  - 100 million users != hundred servers
When considering MTAs, there are typically five things to look at: Security, Speed, Standards, Monitoring, and Management. Classically, open source sendmail has failed more or less on each point, and alternatives such as postfix have done better.

However, more recently open source sendmail has greatly improved on the first three points, to the point where it is arguable which program is best under what circumstances.

Moreover, the commercial version of sendmail has made great strides in the areas of monitoring and management, feats which no other MTA I know of can begin to approach. We currently run postfix on all our externally visible mail servers, but we’ll be happy to rip all that out and replace it with Sendmail Switch.

Sendmail Advanced Message Server is a very new product, and I believe that it will have “one dot oh” problems like all software typically does. However, I also believe that it will probably be better than anything else available, even though it will still have some rough spots.
Detailed Architecture: Retrieval

Sendmail Advanced Message Server

Slave UMD

Slave UMD

Slave UMD

Slave UMD

Master UMD

LDA/MSA

Sendmail Advanced Message Server

POP3

IMAP

TWRU?

Web

WAP

Fax

Voice

OpenLDAP

iPlanet LDAP?
SMTP/POP3 Benchmarking

- Standard Performance Evaluation Committee
  - SPECmail2001
    <http://www.spec.org/osg/mail2001/>
- Russell Coker
  - postal
- Dan Christian, Mozilla Organization
  - mstone
    <http://www.mozilla.org/projects/mstone/>
SMTP/POP3 Benchmarking

- Wietse Venema
  - smtpsink & smtpstone
    <http://www.postfix.org/>

- Yasushi Saito
  - porctest
    <http://porcupine.cs.washington.edu/porcl/distribution.html>

- Stalker Software
  - SMTPTest & POP3Test
    <http://www.stalker.com/MailTests/>
SMTP/POP3 Benchmarking

- dREIC Systems
  - DeJam Analyzing Suite (Java)
    <http://www.dejam.de/>

- Quest Software
  - Benchmark Factory (NT)
    <http://www.benchmarkfactory.com/benchmark_factory/>

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  17th ACM Symposium on Operating System Principles (SOSP ‘99), December 1999
Questions?

- Slides **will** be made available
  - Via USENIX/SAGE web site
  - Or via my “papers” sub-page
    <http://www.shub-internet.org/brad/papers/>
  - At very least, will be linked from my “papers” sub-page